Linear Codes from the Axiomatic Viewpoint

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0. Overview of lectures

- Basic terminology
- Duality and weight enumerators
- Extension problem



Basic vocabulary

- Let *R* be a finite associative ring with 1.
- Let A be a finite unital left R-module; A will be the alphabet.
- A left *R*-linear code over *A* of length *n* is a left *R*-submodule *C* ⊆ *Aⁿ*.
- Right linear codes are defined similarly.
- ► Due to Nechaev and collaborators, 1999.

Weights

- A weight on A is any function w : A → C with w(0) = 0.
- Extend to $w : A^n \to \mathbb{C}$ by

$$w(a_1, a_2, \ldots, a_n) = \sum_{i=1}^n w(a_i).$$

• Restrict w to linear code $C \subseteq A^n$.

Examples

 Hamming weight: for any alphabet A, define the Hamming weight wt by

$$\operatorname{wt}(a) = egin{cases} 1, & a
eq 0, \ 0, & a = 0. \end{cases}$$

- Lee weight: for R = A = Z/NZ, restrict −N/2 < a ≤ N/2 and set w_L(a) = |a| (ordinary absolute value).
- Homogeneous weight (later).

Hamming weight enumerator

For a linear code C ⊆ Aⁿ, define the Hamming weight enumerator of C by

$$\mathsf{hwe}_{\mathcal{C}}(X,Y) = \sum_{x \in \mathcal{C}} X^{n-\mathsf{wt}(x)} Y^{\mathsf{wt}(x)}.$$

hwe_C(X, Y) = ∑ⁿ_{i=0} A_iXⁿ⁻ⁱYⁱ, where A_i is the number of codewords in C of Hamming weight i.



Dual codes

• Let A = R itself. Define the **dot product** on R^n by

$$x \cdot y = \sum_{i=1}^n x_i y_i,$$

where $x = (x_1, x_2, ..., x_n)$ and $y = (y_1, y_2, ..., y_n)$. For a linear code $C \subseteq R^n$, define **annihilators**

$$I(C) = \{ y \in R^n : y \cdot C = 0 \},\$$

$$r(C) = \{ y \in R^n : C \cdot y = 0 \}.$$

Questions

- Are the annihilators well-behaved?
- Is there a nice relationship between the Hamming weight enumerators of C and its annihilators?
 - Yes: the MacWilliams identities.
- ▶ We will discuss these questions later today.



Isometries

- Let $C_1, C_2 \subseteq A^n$ be two linear codes. An *R*-linear isomorphism $f : C_1 \rightarrow C_2$ is a linear **isometry** with respect to a weight *w* if w(xf) = w(x) for all $x \in C_1$.
- I will usually write homomorphisms of left modules M on the right side, so that (rx)f = r(xf) for r ∈ R, x ∈ M.

Symmetry groups

Suppose the alphabet A has weight w. Define
 symmetry groups by

$$egin{aligned} G_{\mathsf{lt}} &= \{u \in \mathcal{U}(R) : w(ua) = w(a), a \in A\}, \ G_{\mathsf{rt}} &= \{\phi \in GL_R(A) : w(a\phi) = w(a), a \in A\}. \end{aligned}$$

Here, U(R) is the group of units of R, and GL_R(A) is the group of invertible R-linear homomorphisms of A to itself.

Monomial transformations

Let G ⊆ GL_R(A) be a subgroup. A G-monomial transformation of Aⁿ is an invertible R-linear homomorphism T : Aⁿ → Aⁿ of the form

$$(a_1, a_2, \ldots, a_n)T = (a_{\sigma(1)}\phi_1, a_{\sigma(2)}\phi_2, \ldots, a_{\sigma(n)}\phi_n),$$

for $(a_1, a_2, \ldots, a_n) \in A^n.$
Here, σ is a permutation of $\{1, 2, \ldots, n\}$ and $\phi_i \in G$ for $i = 1, 2, \ldots, n.$

Monomial transformations are isometries

- Easy: G_{rt}-monomial transformations are isometries of Aⁿ with respect to the weight w.
- Let $C_1 \subseteq A^n$ be a linear code and let T be a G_{rt} -monomial transformation of A^n . Set $C_2 = C_1 T$. Then the restriction of T to C_1 is a linear isometry from C_1 to C_2 .
- Is the converse true? That is, does every linear isometry between linear codes extend to a G_{rt}-monomial transformation? Call this the "Extension Problem."
- More on this in the days ahead.

1. Characters

- Definitions
- Properties
- Fourier transform
- Character modules
- Generating characters

Definitions

- Let A be a finite abelian group (additive notation);
 A will be a module later.
- A character of A is a group homomorphism

$$\pi: \mathbf{A} \to \mathbb{C}^{\times},$$

where \mathbb{C}^{\times} is the multiplicative group of nonzero complex numbers.

- The set of all characters of A is a multiplicative abelian group under pointwise multiplication.
- Additive version: $\widehat{A} \cong \operatorname{Hom}_{\mathbb{Z}}(A, \mathbb{Q}/\mathbb{Z}).$

Duality functor

- Pontryagin duality: $A \mapsto \widehat{A}$
- $\widehat{\widehat{A}} \cong A$, naturally.
- $\widehat{A} \cong A$, but not naturally.
- $\blacktriangleright (A \times B)^{\widehat{}} \cong \widehat{A} \times \widehat{B}.$
- Exact contravariant functor:

$$0 \to A_1 \to A_2 \to A_3 \to 0$$

induces

$$0 \rightarrow \widehat{A}_3 \rightarrow \widehat{A}_2 \rightarrow \widehat{A}_1 \rightarrow 0.$$

Summation formulas

For $a \in A$,

$$\sum_{\pi\in\widehat{A}}\pi(a)=egin{cases} |A|,&a=0,\ 0,&a
eq 0. \end{cases}$$

For
$$\pi \in \widehat{A}$$
,

$$\sum_{\boldsymbol{a}\in A} \pi(\boldsymbol{a}) = \begin{cases} |\boldsymbol{A}|, & \pi = 1, \\ 0, & \pi \neq 1. \end{cases}$$

Linear independence

- Let F(A, C) = {f : A → C}, a complex vector space of dimension |A|.
- The elements of \widehat{A} form a basis for $F(A, \mathbb{C})$.
- ► In particular, characters are linearly independent over C.
- Need multiplicative form of characters for linear independence and for the summation formulas.



Annihilators

Let B ⊆ A be any subgroup.
Define the annihilator (Â : B):

$$(\widehat{\mathsf{A}}:\mathsf{B})=\{\pi\in\widehat{\mathsf{A}}:\pi(\mathsf{B})=1\}.$$

- $(\widehat{A}:B)\cong (A/B)^{\widehat{}}.$
- $|B| \cdot |(\widehat{A} : B)| = |A|.$
- Double annihilator: $(A : (\widehat{A} : B)) = B$.



Fourier transform

• Given a function $f : A \to V$, V a complex vector space. Define its **Fourier transform** $\hat{f} : \hat{A} \to V$ by

$$\widehat{f}(\pi) = \sum_{a \in \mathcal{A}} \pi(a) f(a), \quad \pi \in \widehat{\mathcal{A}}.$$

$$\widehat{}: F(A, V) \to F(\widehat{A}, V).$$

Invert:

$$f(a)=rac{1}{|\mathcal{A}|}\sum_{\pi\in\widehat{\mathcal{A}}}\pi(-a)\widehat{f}(\pi), \quad a\in\mathcal{A}.$$

Poisson summation formula

Let *B* be any subgroup of *A*, and let $f : A \rightarrow V$. Then for any $a \in A$,

$$\sum_{b\in B}f(a+b)=rac{1}{|(\widehat{A}:B)|}\sum_{\pi\in(\widehat{A}:B)}\pi(-a)\widehat{f}(\pi).$$

If a = 0, then

$$\sum_{b\in B} f(b) = \frac{1}{|(\widehat{A}:B)|} \sum_{\pi\in(\widehat{A}:B)} \widehat{f}(\pi).$$

Character modules

- Now suppose R is a finite ring with 1 and A is a finite unital left R-module.
- Then \widehat{A} becomes a right *R*-module by

$$\pi^r(a) = \pi(ra), \quad a \in A.$$

- $({}^{r}\pi(a) = \pi(ar)$ for right to left case.)
- ▶ ˆ is an exact contravariant functor of *R*-modules.
- For left *R*-submodule *B* ⊆ *A*, (Â : *B*) is a right *R*-submodule of Â.

Top-bottom duality

- An *R*-module is **simple** if it has no nontrivial proper submodules.
- The Jacobson radical Rad(R) is the intersection of all maximal left ideals of R.
- For a left *R*-module *A*, the socle Soc(*A*) is the left *R*-submodule generated by all simple left *R*-submodules of *A*.

•
$$(A/\operatorname{Rad}(A)A)^{\widehat{}} \cong \operatorname{Soc}(\widehat{A})$$

Generating characters

- Left *R*-module *A*.
- A character ρ ∈ Â is a generating character if ker ρ contains no nonzero left R-submodules.
- ► Not every module admits a generating character.



Embedding

- Suppose ρ is a generating character for A.
- Define $\alpha : A \to \widehat{R}$ by $(a\alpha)(r) = \rho(ra)$, $a \in A$, $r \in R$.
- α is an injective homomorphism of left *R*-modules.
- Dual map $R \to \widehat{A}$, $r \mapsto \rho^r$, is surjective homomorphism of right *R*-modules.
- ρ generates \widehat{A} .
- Conversely: if is cyclic, or A embeds in R, then A has a generating character.

Frobenius rings

- Recall: $|\widehat{R}| = |R|$.
- Consider *R* as a module over itself.
- If R has a generating character, then R
 [^] ≅ R as left and as right R-modules.
- $\operatorname{Soc}(R) = \operatorname{Soc}(\widehat{R}) \cong (R/\operatorname{Rad}(R))^{\widehat{}} \cong R/\operatorname{Rad}(R).$
- Such a finite ring is a Frobenius ring.

Some examples of generating characters

•
$$\mathbb{Z}/N\mathbb{Z}$$
 admits $\theta_N(a) = \exp(2\pi i a/N)$, $a \in \mathbb{Z}/N\mathbb{Z}$.

- exp is the standard complex exponential function.
- \mathbb{F}_q admits $\theta_q(a) = \theta_p(\operatorname{Tr}_{q \to p}(a)), \ a \in \mathbb{F}_q.$
- $R = M_{k \times k}(\mathbb{F}_q)$ admits $\rho(P) = \theta_q(\operatorname{Tr} P)$, $P \in R$.
- $A = M_{k imes \ell}(\mathbb{F}_q)$, $k > \ell$, admits $\rho|_A$.
- When k < ℓ, A does not admit a generating character. π_Q(P) = θ_q(Tr(PQ)), for Q ∈ M_{ℓ×k}(𝔽_q). Find nonzero X ∈ M_{k×ℓ}(𝔽_q) with XQ = 0, as k < ℓ. Then RX ⊆ ker π_Q.



Cyclic socle

- ► Suppose *A* has cyclic socle Soc(*A*).
- Soc(A) is a sum of matrix modules with $k \ge \ell$.
- Soc(A) admits a generating character: multiply together those from matrix modules.
- Extension exists, by exactness. Any extension is a generating character for A.